

# Continuous Time Quantum Walk on finite dimensions

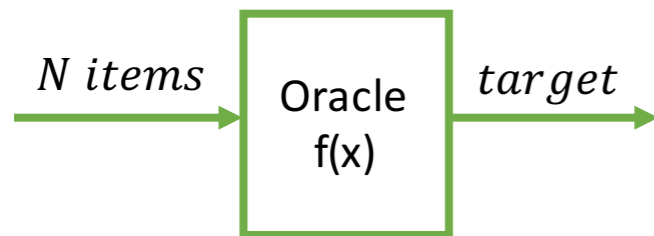
*Shanshan Li*

*Joint work with Stefan Boettcher*

*Emory University*

*QMath13, 10/11/2016*

# Grover Algorithm: Unstructured Search



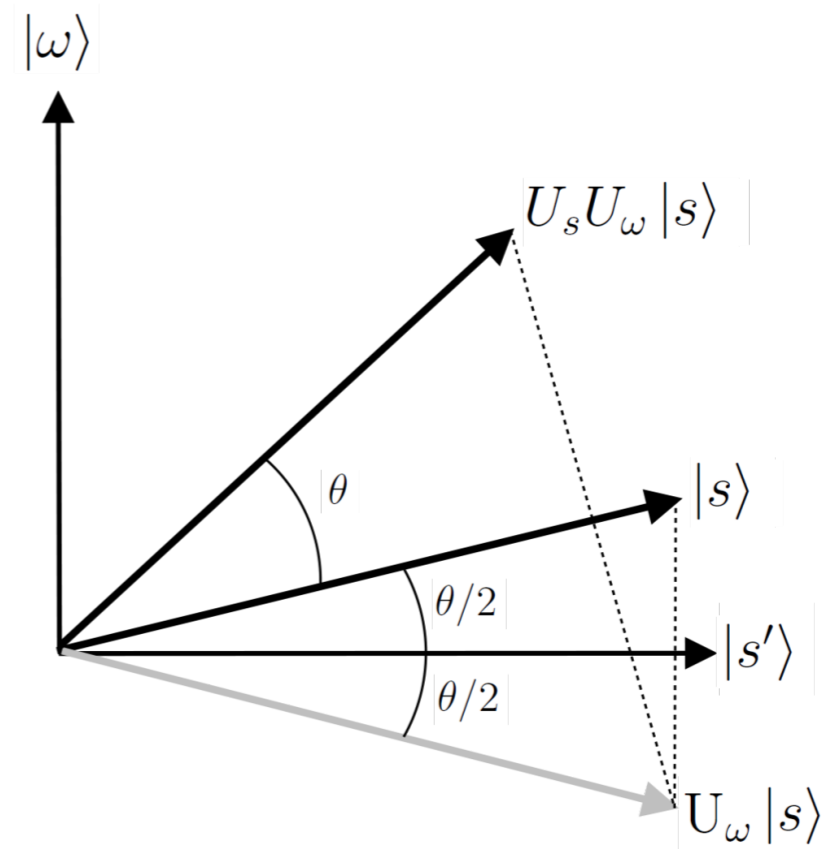
$$f(x) = \begin{cases} 1 & x = w \\ 0 & \text{otherwise} \end{cases}$$

Initialize the system to the state

$$|s\rangle = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle$$

Apply Grover Iteration

$$(U_s U_w)^{\sqrt{N}} |s\rangle$$



$$U_s = 2 |s\rangle \langle s| - I$$

$$U_w = I - 2 |w\rangle \langle w|$$

$$\theta = 2 \arcsin \frac{1}{\sqrt{N}}$$

# Quantum Walk Basics for Spatial Search

Random Walk

$$\frac{d}{dt}p_x = \sum_{y=0}^{N-1} L_{xy}p_y$$

Continuous time quantum walk

$$\frac{d\Psi_x(t)}{dt} = \sum_y H_{xy}\Psi_y(t)$$

$$H = \gamma L - |\omega\rangle\langle\omega|$$

$$\text{initial state } |s\rangle = \frac{1}{N} \sum_x |x\rangle$$

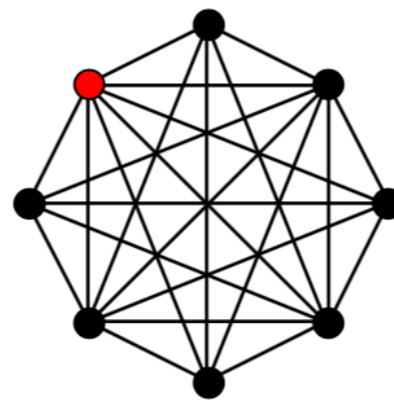
marked state  $|\omega\rangle$

**Unstructured search:**

$f(x)$  is a computable function

**Spatial Search:**

N items stored in a d-dimensional physical space

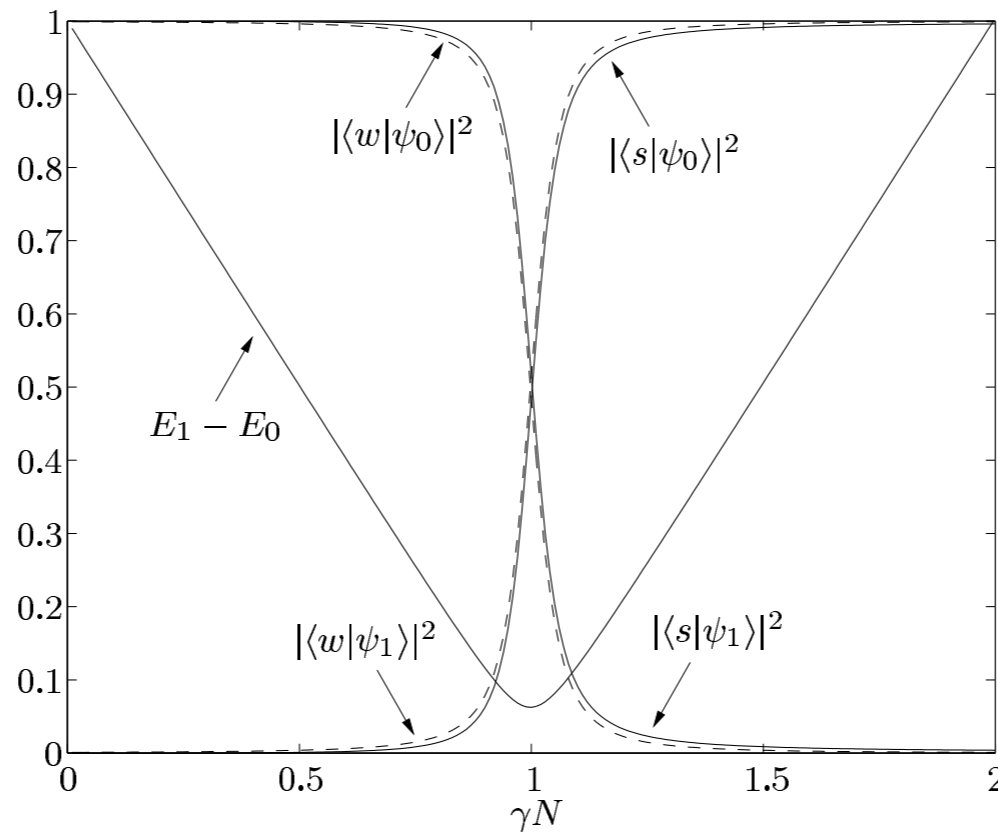


Laplacian = Degree Matrix - Adjacency Matrix

$$L = D - A$$

# Critical Point in the Hamiltonian

$$\begin{aligned} \gamma &= 0 \\ H &= -|w\rangle\langle w| \\ 0^{th}, 1^{st} &= |w\rangle, |s\rangle \end{aligned}$$



$$\begin{aligned} \gamma &= \infty \\ H &= \gamma L \\ 0^{th}, 1^{st} &= |s\rangle, |w\rangle \end{aligned}$$

$$\begin{aligned} \gamma &= \gamma_c \\ 0^{th}, 1^{st} &= (|w\rangle \pm |s\rangle) / \sqrt{2} \\ T &= \frac{\pi}{2} \sqrt{N} \end{aligned}$$

---

# CTQW: optimal performance

---

Quadratic Speedup  $O(\sqrt{N})$  **Grover efficiency**

- complete graph, hypercube, strongly regular graph

(E. Farhi and S. Gutmann 1998, A. M. Childs et al 2002, J. Janmark et al, 2014)

- Erdős Renyi graph  $p \geq \log^{\frac{3}{2}} N/N$

(S. Chakraborty et al, 2016)

- lattices  $d > 4$

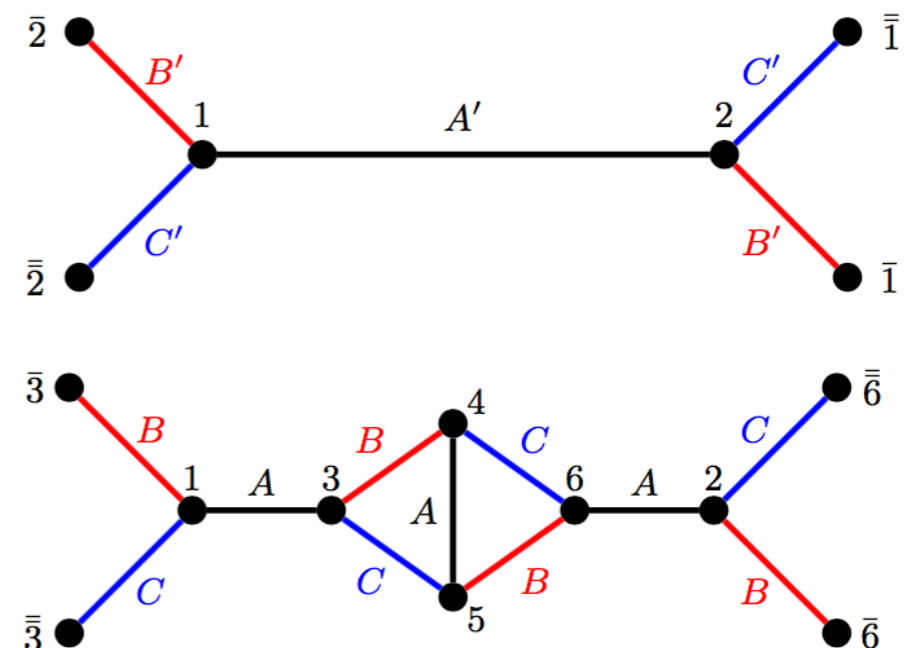
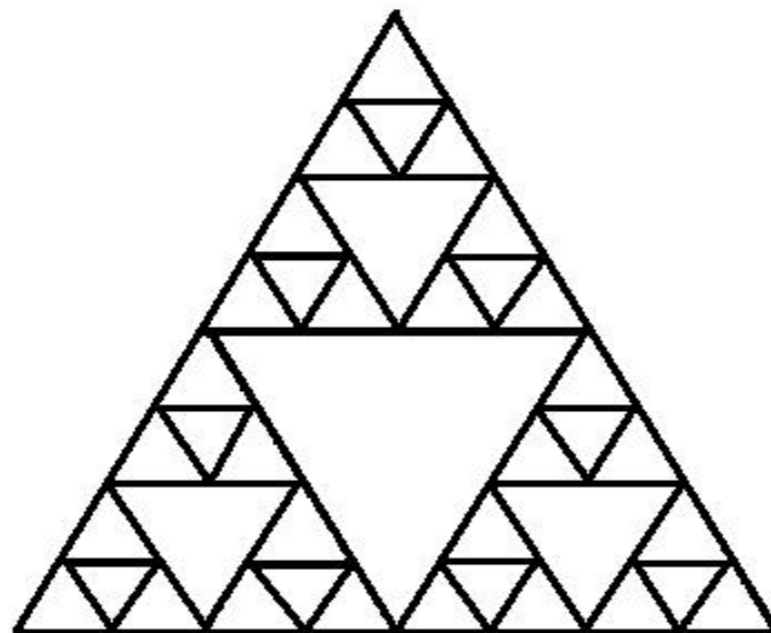
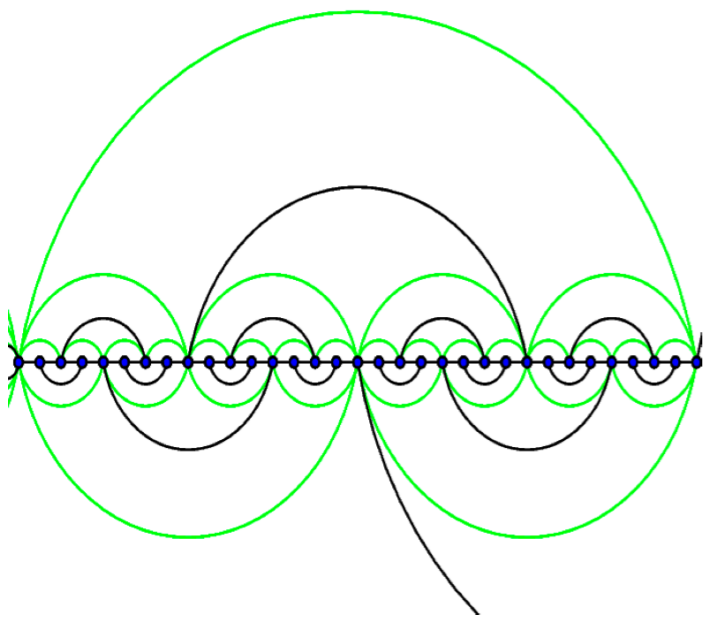
(A. M. Childs et al 2004)

- ❖ We extend CTQW to fractal graphs with real fractal dimension
- ❖ Spectral dimension of graph Laplacian determines the computational complexity

# Finite Dimensional Fractals

we generalize to arbitrary real (fractal) dimension

- ❖ Hierarchical networks
- ❖ Sierpinski Gasket, Migdal-Kadanoff network with regular degree 3
- ❖ **Diamond fractals based on the Migdal-Kadanoff renormalization group scheme**



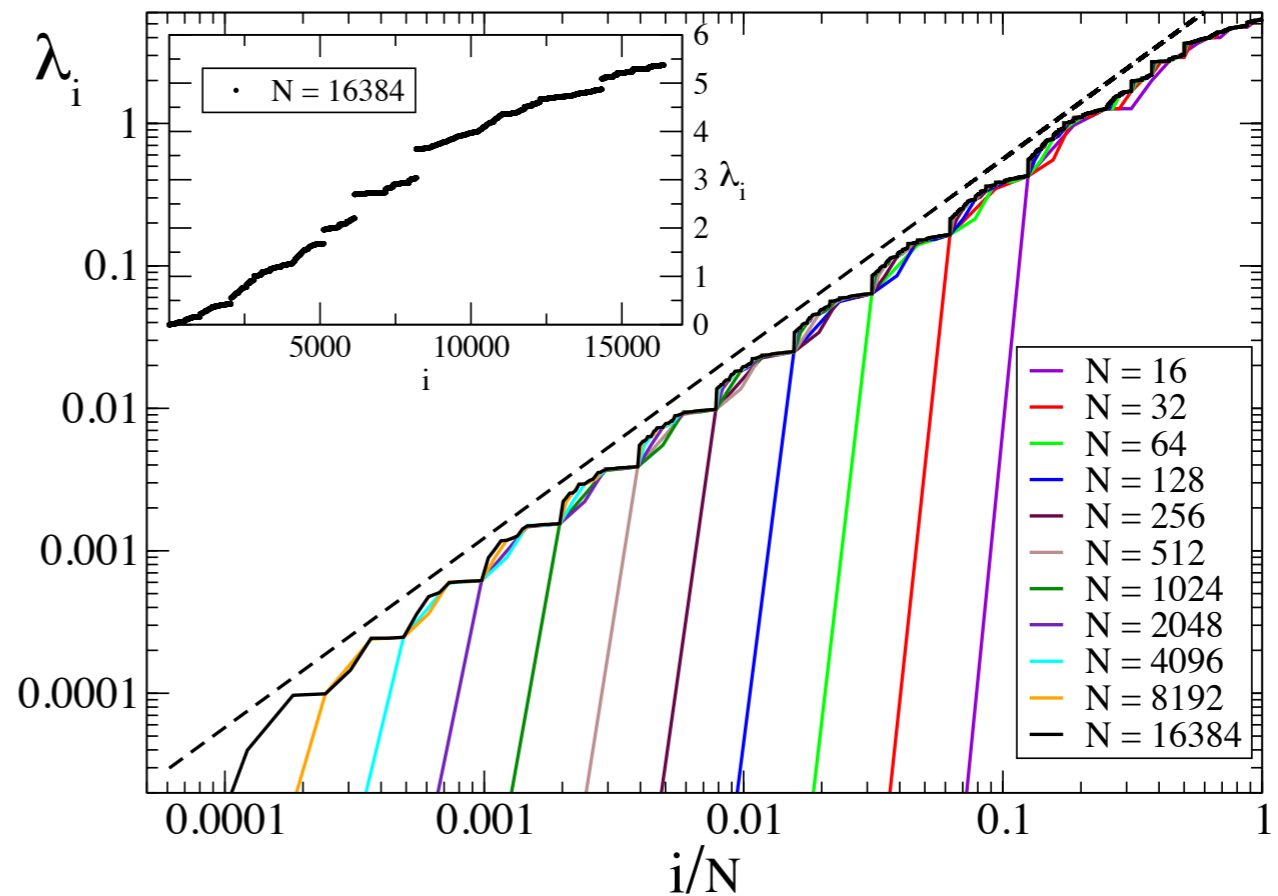
# Dimensions in fractal networks

Fractal dimensions

$$N \sim l^{d_f}$$

Spectral dimension

$$\lambda_i \sim N^{-2/d_s}$$

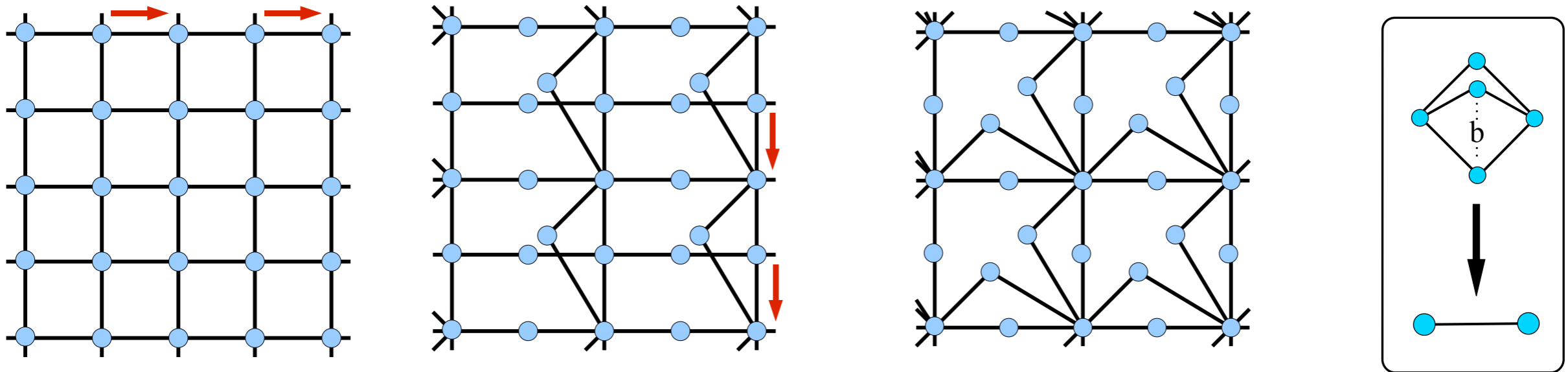


Hierarchical Network with regular degree 3

# Migdal-Kadanoff renormalization group(MKRG)

Model regular lattices closely  $d_s = d_f = d$

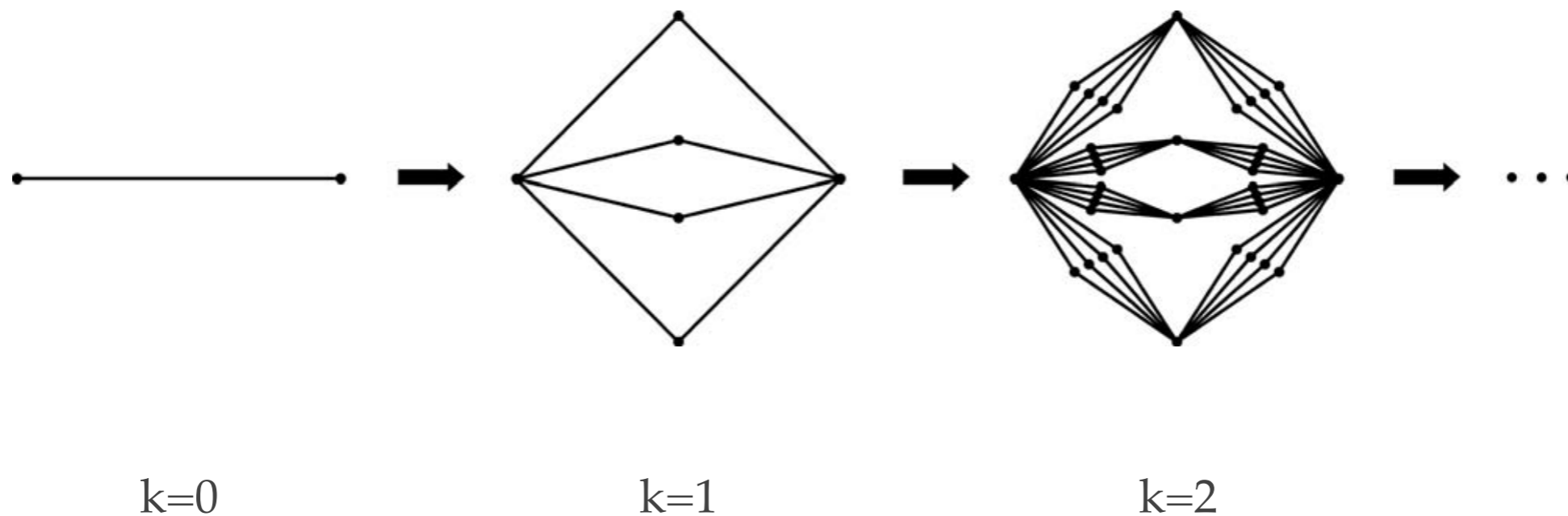
arbitrary real dimension



Bond-moving scheme on square lattices with rescaling length  $l=2$ ,  
branching factor  $b=2$



# Procedure to build the Diamond Fractals



$$d_f = d_s = d = 1 + \frac{\ln b}{\ln 2}$$

---

# Measure Critical Point

---

The spectral Zeta function

$$I_j = \frac{1}{N} \sum_{i=1}^{N-1} \left(\frac{1}{\lambda_i}\right)^j$$

When the CTQW is optimal for search, the critical point takes place (numerically true for almost all sites in fractals we consider)

$$\gamma_c = I_1$$

The transition probability

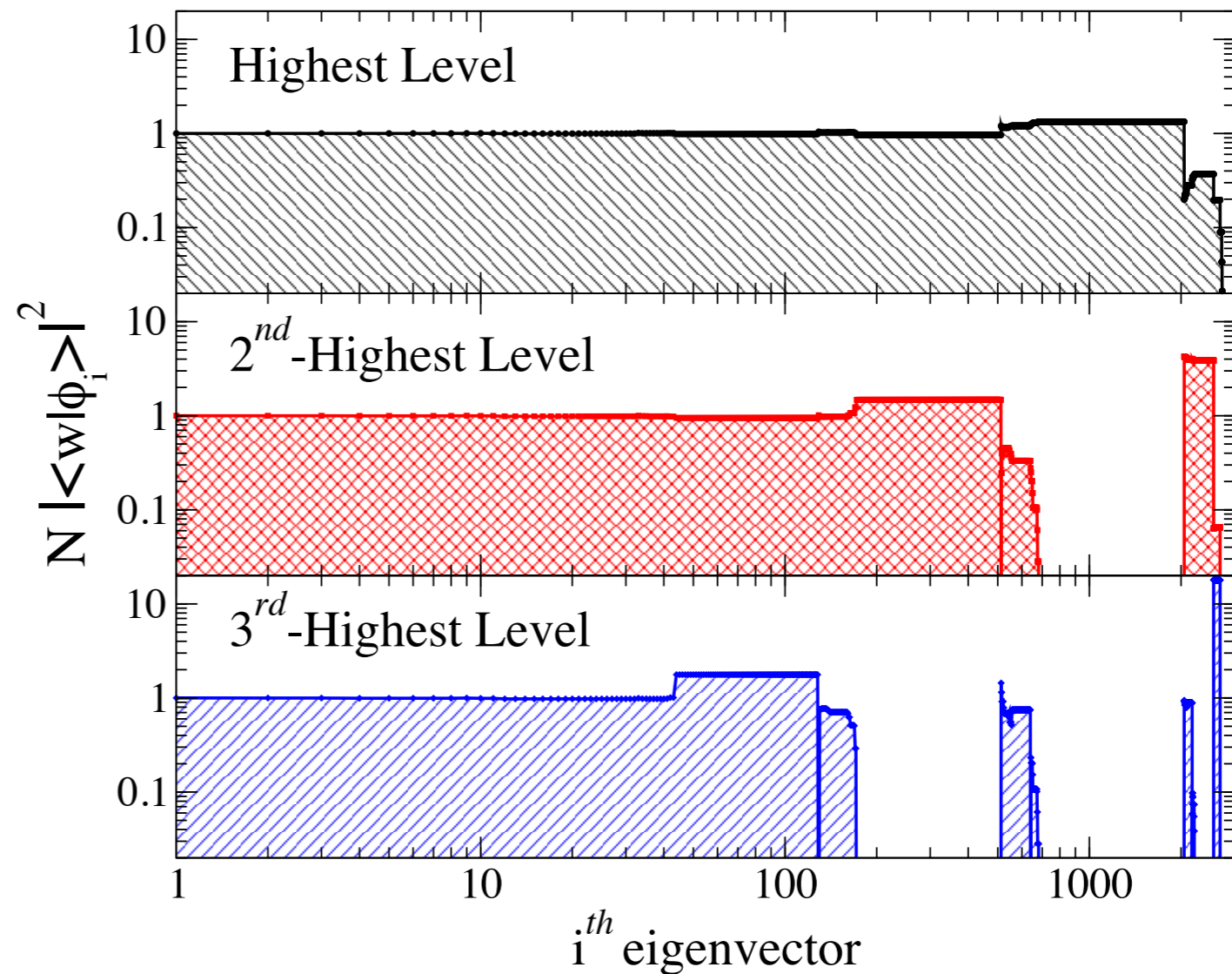
$$\Pi_{s,\omega} = \frac{I_1^2}{4I_2} \sin^2\left(\frac{2I_1}{\sqrt{I_2}} \frac{t}{\sqrt{N}}\right)$$

# Assumption on fractal Laplacian eigenvector

$$I_j = \sum_{i=1}^{N-1} \frac{|\langle \omega | \phi_i \rangle|^2}{\lambda_i^j}$$

$$|\langle \omega | \phi_i \rangle|^2 \sim \frac{1}{N}$$

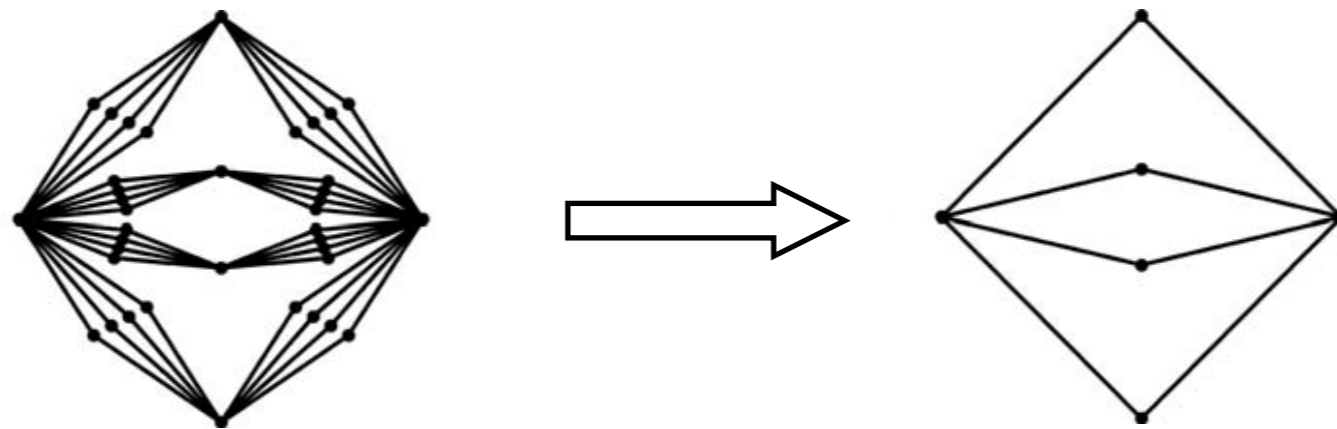
MK renormalization group with  $b=2$



---

# Renormalization Group Argument

---



$$\det \left[ L^{(k)} \left( q_i^{(k)}, p_i^{(k)}, \dots \right) \right]$$

$$\det \left[ L^{(k+1)} \left( q_i^{(k+1)}, p_i^{(k+1)}, \dots \right) \right]$$

The spectral Zeta function

$$\begin{aligned} I_j &\sim \left( \frac{\partial}{\partial \epsilon} \right)^j \ln \left[ \frac{1}{\epsilon} \det (L + \epsilon) \right] \Big|_{\epsilon \rightarrow 0} \\ &\sim \begin{cases} N^{\frac{2j}{d_s} - 1} & d_s < 2j \\ \text{const} & d_s > 2j \end{cases} \end{aligned}$$

---

# Computational Complexity of CTQW

---

$$\gamma_c \sim \begin{cases} N^{\frac{2}{d_s}-1}, & d_s < 2 \\ \text{const}, & d_s > 2 \end{cases} \quad T \sim \begin{cases} N^{\frac{1}{2}}, & d_s > 4 \\ N^{\frac{1}{2}} \ln^{\frac{3}{2}} N, & d_s > 2 \\ \gtrsim N^{\frac{2}{d_s}}, & 2 < d_s < 4 \end{cases}$$

- ❖ spectral dimension of network Laplacian determines the computational complexity
- ❖ complement the discussions on regular lattices and mean-field networks

## References:

- ❖ Shanshan Li and Stefan Boettcher, arXiv preprint arXiv: 1607.05317 , 2016
- ❖ Stefan Boettcher and Shanshan Li, arXiv preprint arXiv: 1607.05168 , 2016

**Thank you !**

# RG calculation for spectral determinant

$$\frac{1}{\sqrt{\det L}} = \int \cdots \int_{-\infty}^{\infty} \prod_{i=1}^N \frac{dx_i}{\sqrt{\pi}} \exp \left\{ - \sum_{n=1}^N \sum_{m=1}^N x_n L_{n,m} x_m \right\}$$

$$B_i(x, y) = C_i \exp \left\{ -\frac{q_i}{2} x^2 - \frac{q_0}{2} y^2 + 2pxy \right\} \Rightarrow$$

$$B'_{i-1}(x, z) = \int \cdots \int_{-\infty}^{\infty} \prod_{j=1}^b \frac{dy_j}{\sqrt{\pi}} B_i(x, y_j) B_1(y_j, z) \Rightarrow$$

$$B'_{i-1}(x, z) = C'_{i-1} \exp \left\{ -\frac{q'_{i-1}}{2} x^2 - \frac{q'_0}{2} y^2 + 2p'xz \right\}$$

$$\det(L + \epsilon)|_{\epsilon \rightarrow 0} = C_k^{-2} b^{2k} (q_k^2/4 - p_k^2)$$

